

## IN THE SPECIFICATION

Please replace the paragraph beginning at page 8, line 15 with the following:

--The lock modes described above are superior to other lock modes. Assume two non-serializable operations delete separate rows B and C such that row B is next to row C, and row C is next to row D in a non-unique index. In other DBMSs, the delete for row B takes an X lock on row B and a X (or other proprietary weaker locking modes e.g. NX in IBM's DB2) lock on row C. At this point the delete for row C cannot proceed since it needs a X lock (or at least a W lock in IBM's DB2), on row C. In the neighborhood locking scheme, the delete for row B can take a Xn ~~Xw~~ lock on row B and a Xnei lock on row C. The delete for row C takes a Xn ~~Xw~~ lock on row C, which is compatible and an Xnei lock on row C. Thus both deletes are concurrent.--